

**CSE 451: Operating Systems**  
**Winter 2025**

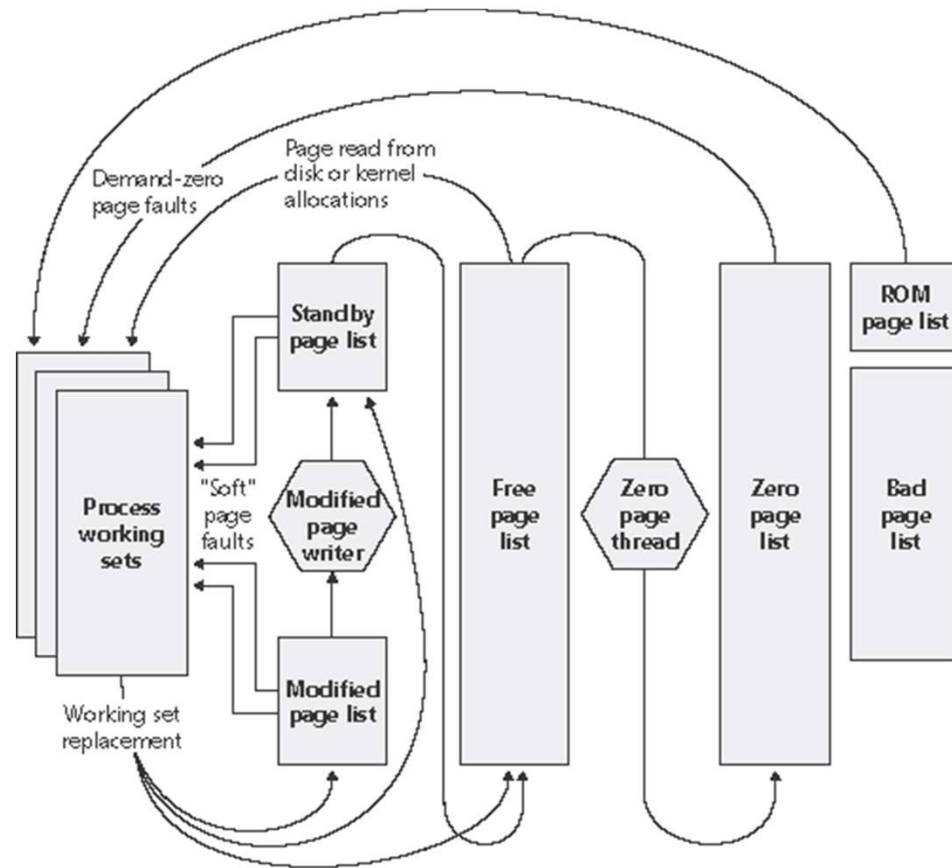
**Module 14**  
**Windows MM**

**Gary Kimura**

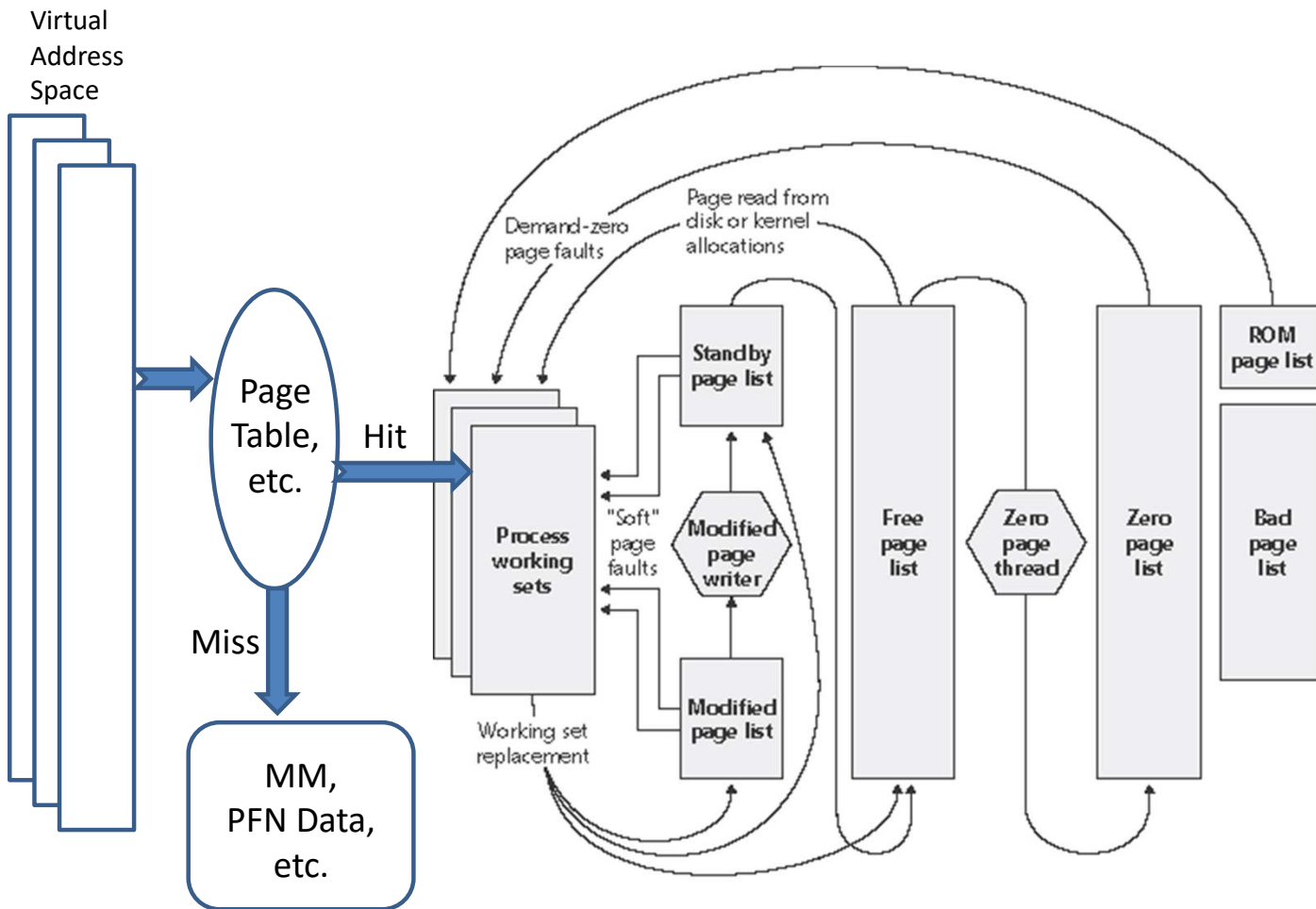
## Memory Management (continued)

- Windows Paging
- Windows Kernel Heap
- Wickedly Fun Exam Question

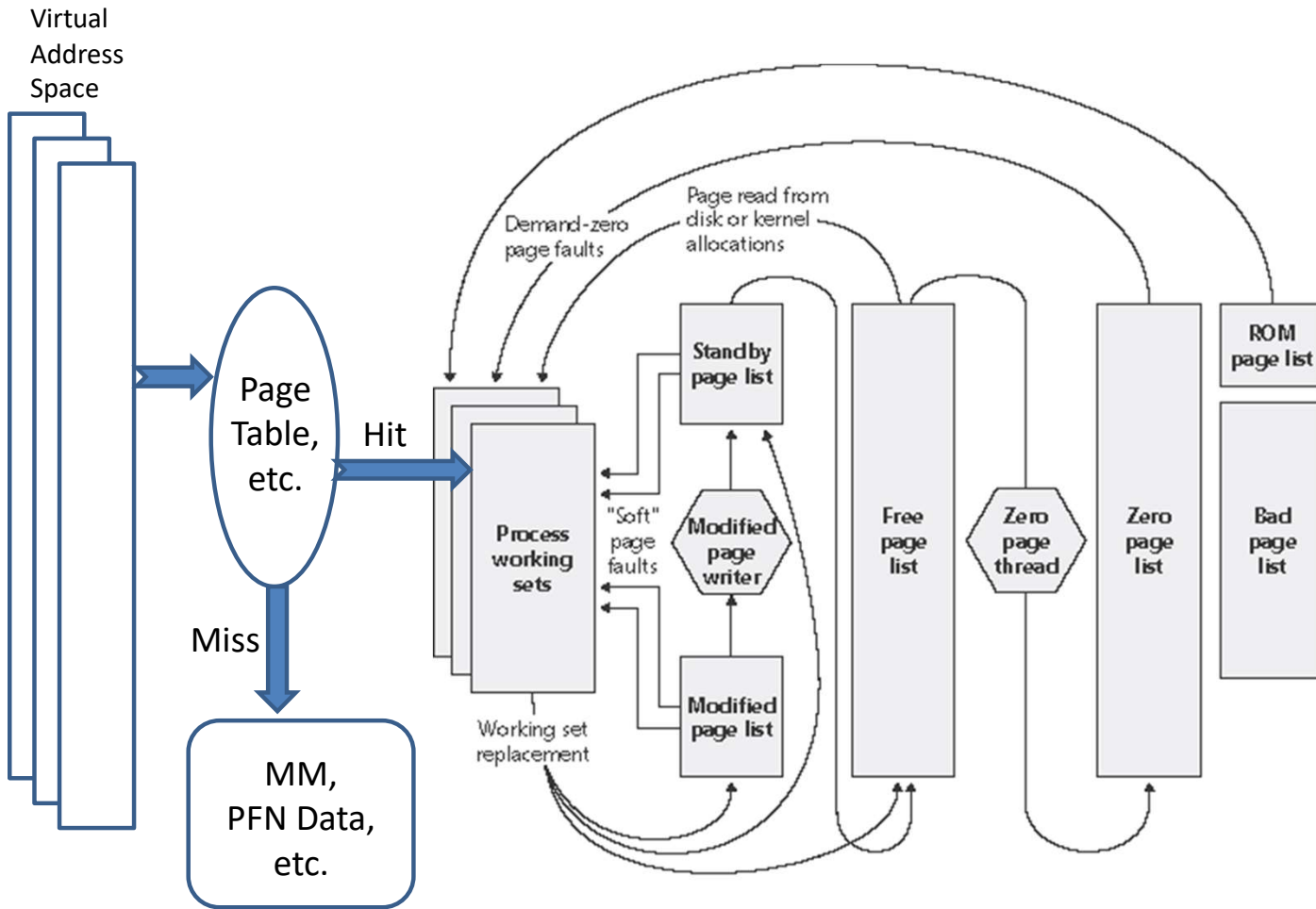
# Windows Page Frame State Diagram



# Windows Page Frame State Diagram



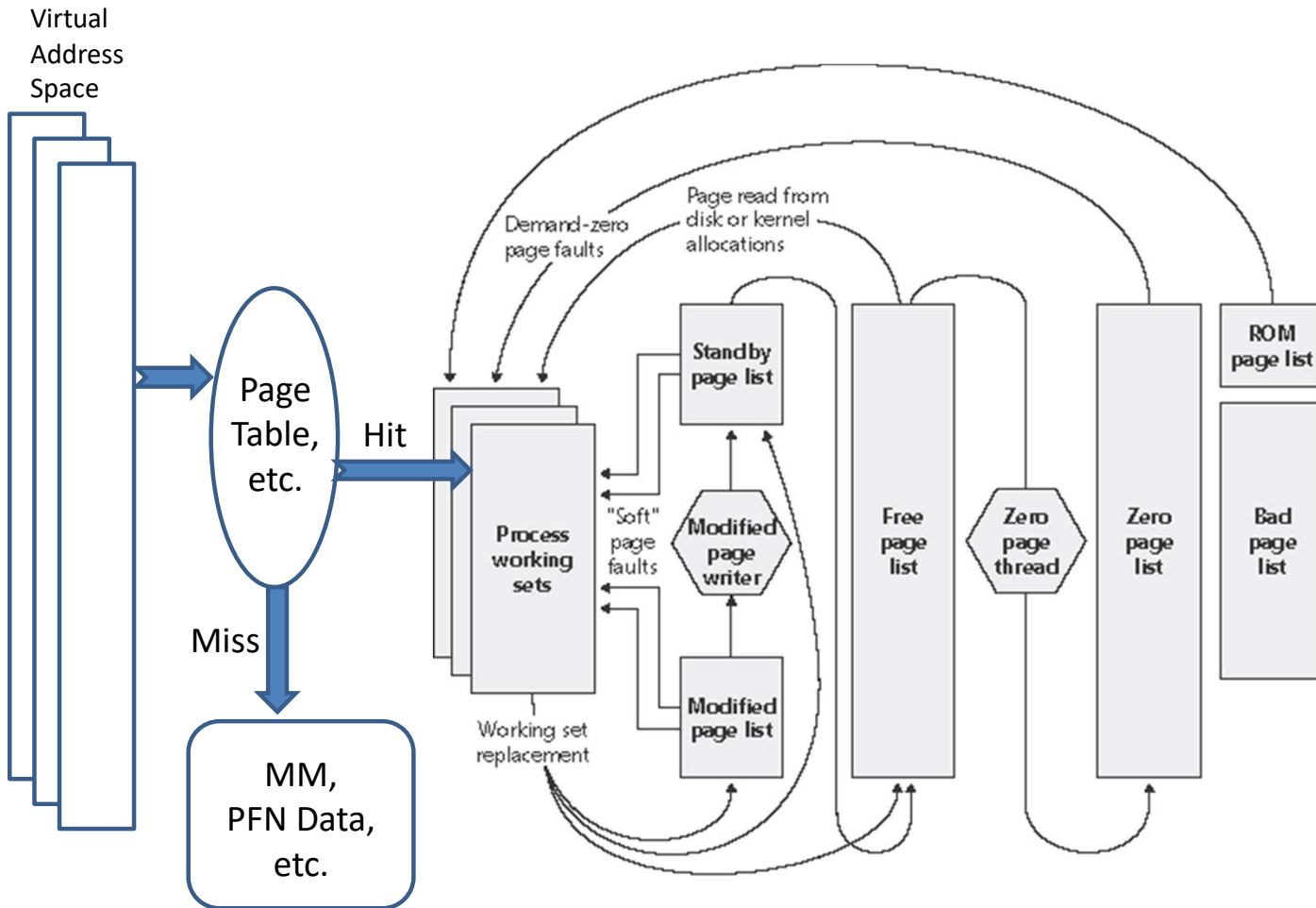
# Windows Page Frame State Diagram



## Page States

- Active (also called Valid)
- Transition
- Standby
- Modified
- Modified no-write
- Free
- Zeroed
- Rom
- Bad

# Windows Page Frame State Diagram



## Paging Features

- Local and Global page replacement
- LRU on top of FIFO
- Hard and Soft page faults

## Windows Kernel Pool (aka Heap)

- Boundary tagged (tried but rejected Fibonacci and buddy system)
- Paged and Nonpaged (Once had Must Succeed)
- Lookaside lists
- Node type codes to help quickly identify objects in the pool

# Overall Pool Layout



## Debugging pool corruption

- Checked build versus Free build
- Debugging pool corruption bugs, often stale pointers or allocation overruns
- 0xDEADBEEF and 0xBAADF00D
- Extra code to check for pool corruption
- A pointer hack I used to catch a bad actor (data alignment fault)

# Pool Corruption

## A Fun Exam Question from 2013

- Examine how long it takes a user mode program writing to an array of integers.
- Assumptions
  - The entire array will fit into physical memory (no paging)
  - The system is pretty much idle except for this program
- First malloc() the array, and then
- Time how long it takes to write to every element of the array using various access patterns.

## The Actual Exam Question

Consider the following program that allocates a multi-megabyte sized array of unsigned longs, and then times how long it takes to write to every array location. The program varies the pattern it uses to write to each array location based on a stride that changes between each pass through the array.

For example a stride of 1 makes one pass through the array accessing locations 0,1,2,... until the end of the array is reached. A stride of 2 makes two passes through the array, first accessing locations 0,2,4,... and then accessing locations 1,3,5,... until the end of the array is reached. The program starts with a stride value of 1 and then increases it, based on user input, until the stride is equal to half the size of the array. The program times how long it takes, in seconds, for each new stride through the array.

The program takes three parameters, first is the number of megabytes to allocate to the array, the second and third parameters are the multiplication and additive factors used to compute the stride. For example, the parameters “2 1 1” allocate a 2MB array testing stride values of 1,2,3,4,..., 131072. Note that 131072 is the halfway point in a 2MB integer array. The parameters “2 2 0” allocate a 2MB array testing stride values of 1,2,4,8,16,...,131072. In other words the stride value doubles each time.

```

void main (int argc, char *argv[])
{
    clock_t StartTime, EndTime;
    unsigned long *Array, Size, StrideTimes, StridePlus, i,j,k;

    sscanf(argv[1], "%lu", &Size);
    sscanf(argv[2], "%lu", &StrideTimes);
    sscanf(argv[4], "%lu", &StridePlus);

    printf("Size = %luMB\n", Size);

    // Allocate a test array
    Size = 1024*1024*Size;
    if ((Array = malloc(Size)) == NULL) {
        printf("malloc failed\n");
        return;
    }
    Size /= 4;

    // Now test it for strides from 1 to size/2
    printf(" Stride  Seconds\n");
    for (i = 1; i < Size/2; i = (i*StrideTimes)+StridePlus) {
        printf("%8lu", i);
        StartTime = clock();
        for (j = 0; j < i; j++) {
            for (k = j; k < Size; k += i) {
                Array[k] = k;
            }
        }
        EndTime = clock();
        printf(", %8.3f\n", ((double)(EndTime - StartTime)/CLOCKS_PER_SEC));
    }
}

```

What about calling free()  
when the program exits?



This program was run on both Windows and Linux systems with 4GB of RAM. Here is the data for a run of “1024 2 0” on a Linux system.

Size = 1024 MB

Stride	Seconds	Stride	Seconds	Stride	Seconds
1,	1.920	512,	11.890	262144,	18.250
2,	1.560	1024,	12.190	524288,	14.710
4,	2.780	2048,	12.960	1048576,	9.810
8,	5.530	4096,	14.280	2097152,	5.310
16,	11.450	8192,	16.510	4194304,	4.100
32,	16.470	16384,	22.070	8388608,	3.820
64,	15.000	32768,	22.390	16777216,	3.760
128,	13.390	65536,	21.510	33554432,	2.170
256,	12.310	131072,	20.270	67108864,	1.170

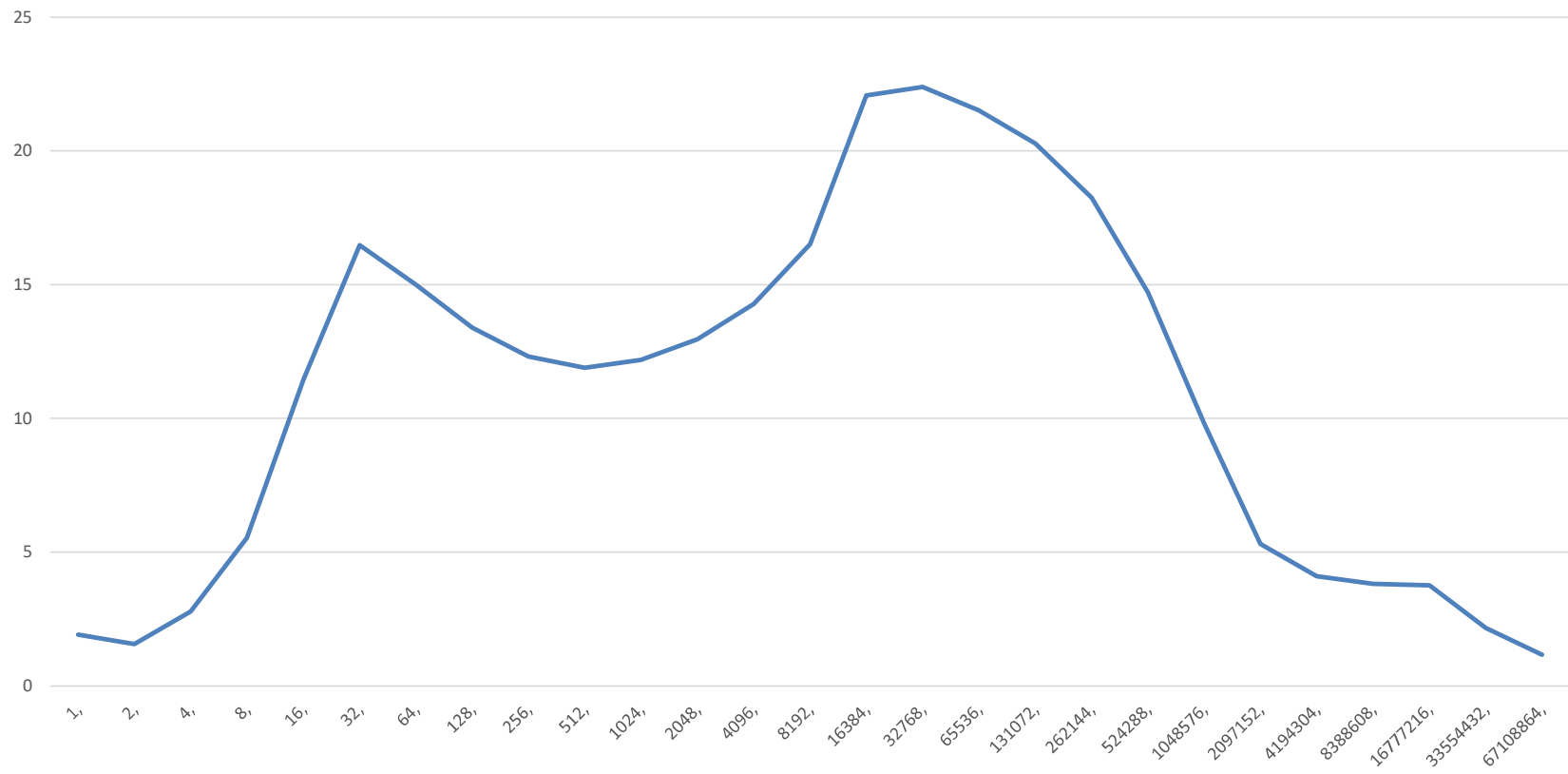
## Two questions to answer

[20 points] Notice how the first pass with a stride of 1 takes longer than the second pass with a stride of 2. This behavior showed up consistently on Linux but not Windows. Please give a plausible explanation for what causes this phenomenon (it might be a mix of both hardware and software), and what the operating system can do to prevent it. You will need to justify your answer.

[20 points] Also notice how the time for each pass increases and then decreases as the stride values grow from 1 to 67108864. Both Windows and Linux exhibited this behavior. Please give a plausible explanation for this phenomenon (it might be a mix of both hardware and software), and what the operating system can do to prevent it. You will need to justify your answer.



# Illustration of data



## Things to consider

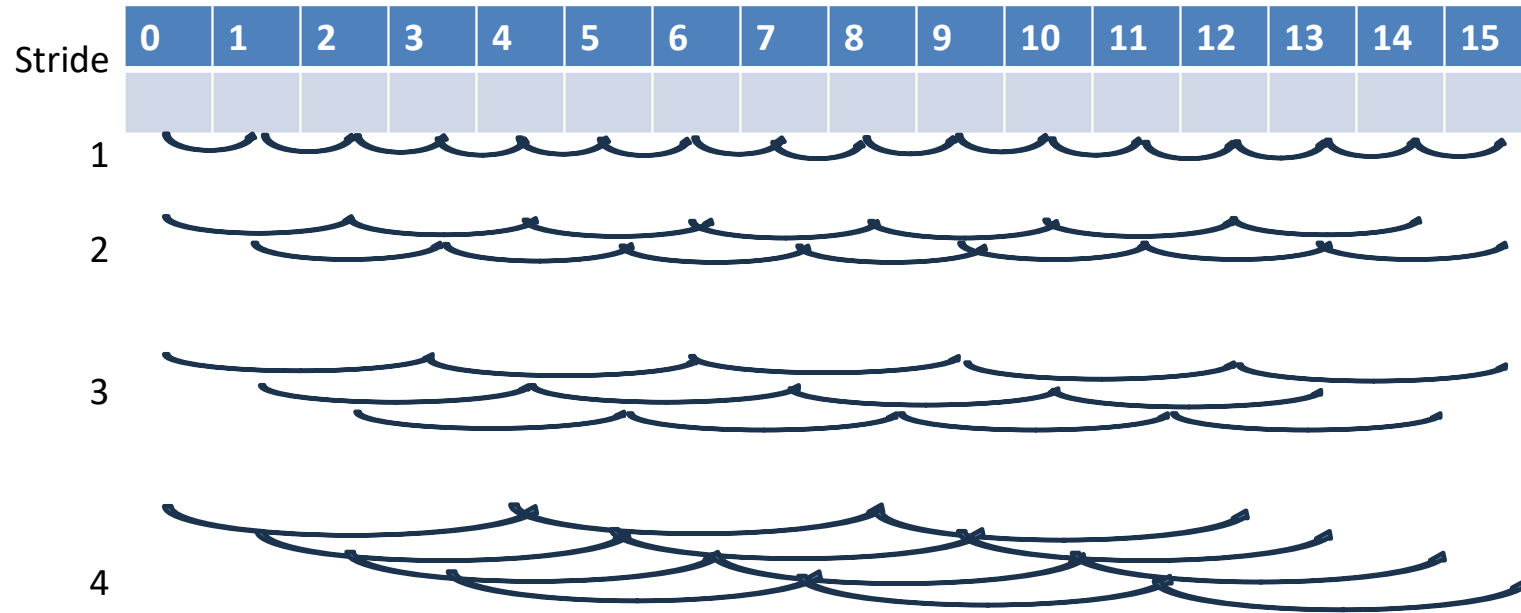
- **Zero pages** – for question 1
- **TLB behavior** – for question 2
- **Various cache levels and Cache line sizes** – for the question not asked about the big dip in the curve

Zero Pages

TLB

## Cache Levels and Cache lines size

# Increasing stride by 1 each time



Now consider the same program run with “64 1 1”. The entire output is too long to include here, but the following is a small section of the output that shows another anomaly that occurs on both Windows and Linux.

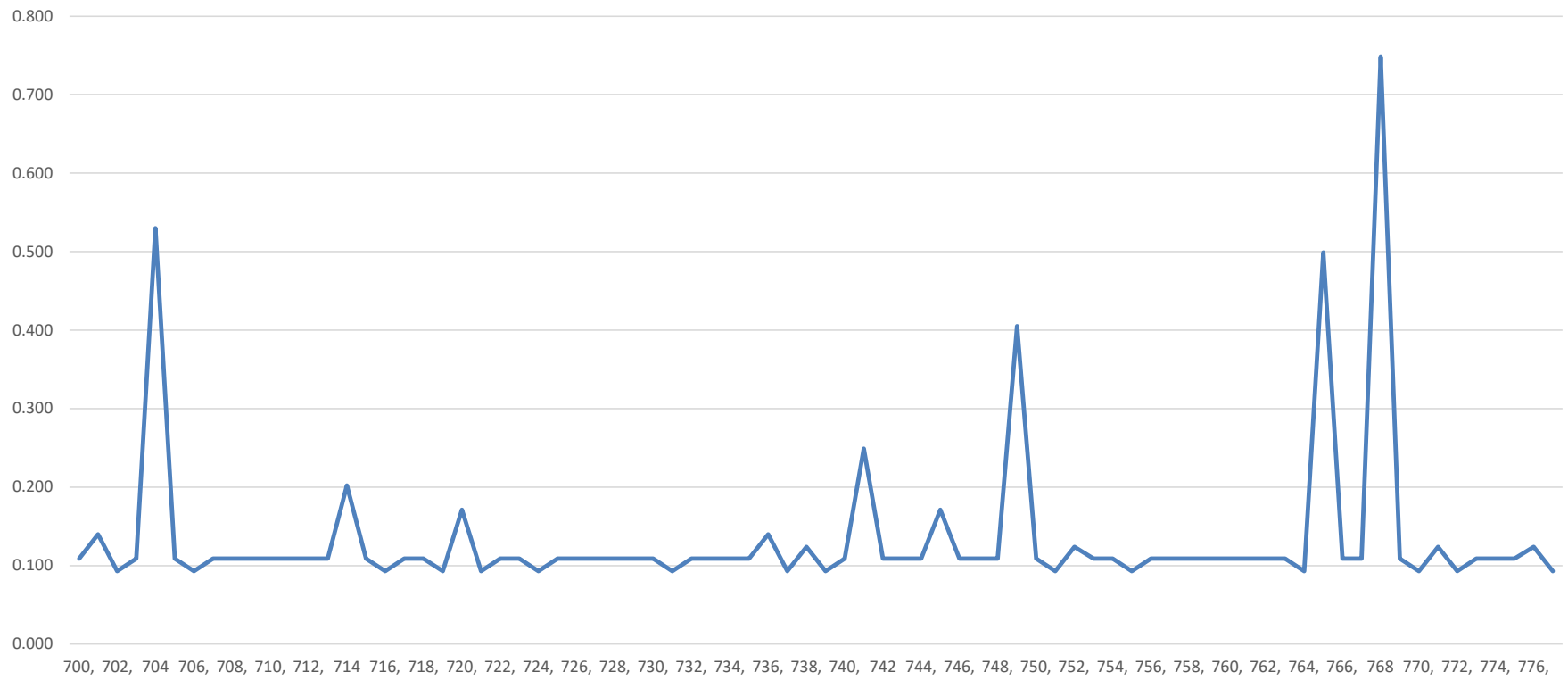
Size = 64MB					
Stride	Seconds	Stride	Seconds	Stride	Seconds
700,	0.109	726,	0.109	752,	0.124
701,	0.140	727,	0.109	753,	0.109
702,	0.093	728,	0.109	754,	0.109
703,	0.109	729,	0.109	755,	0.093
<b><u>704,</u></b>	<b><u>0.530</u></b>	730,	0.109	756,	0.109
705,	0.109	731,	0.093	757,	0.109
706,	0.093	732,	0.109	758,	0.109
707,	0.109	733,	0.109	759,	0.109
708,	0.109	734,	0.109	760,	0.109
709,	0.109	735,	0.109	761,	0.109
710,	0.109	736,	0.140	762,	0.109
711,	0.109	737,	0.093	763,	0.109
712,	0.109	738,	0.124	764,	0.093
713,	0.109	739,	0.093	<b><u>765,</u></b>	<b><u>0.499</u></b>
<b><u>714,</u></b>	<b><u>0.202</u></b>	740,	0.109	766,	0.109
715,	0.109	<b><u>741,</u></b>	<b><u>0.249</u></b>	767,	0.109
716,	0.093	742,	0.109	<b><u>768,</u></b>	<b><u>0.748</u></b>
717,	0.109	743,	0.109	769,	0.109
718,	0.109	744,	0.109	770,	0.093
719,	0.093	745,	0.171	771,	0.124
720,	0.171	746,	0.109	772,	0.093
721,	0.093	747,	0.109	773,	0.109
722,	0.109	748,	0.109	774,	0.109
723,	0.109	<b><u>749,</u></b>	<b><u>0.405</u></b>	775,	0.109
724,	0.093	750,	0.109	776,	0.124
725,	0.109	751,	0.093	777,	0.093

## Question to ponder

[20 points] Notice how the times are consistently in the low 100ms range except for an occasional blip in the 200ms to 700ms range. These blips have been underlined. Please offer an explanation for these blips. Your answer needs to offer a plausible explanation of what is causing this anomaly (it might be a mix of both hardware and software), and what the operating system can do to prevent it, if anything. You will need to justify your answer. If you do not offer an educated guess on what causes this phenomenon explain what you could do to determine its cause.



# Graph of data



## Things to consider

- **Virtual and physical caches** – doesn't answer the question, but good to know
- **Page coloring** – answers the question

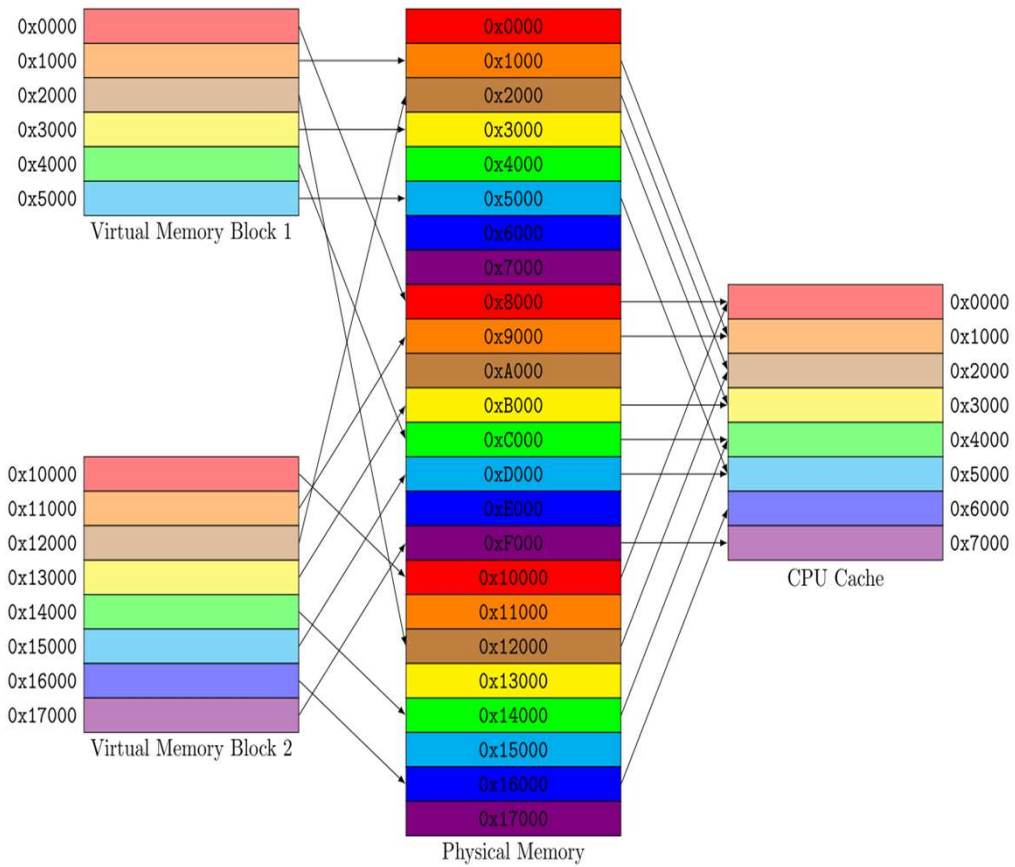
# Virtual and Physical caches

## Virtual and Physical caches

- Is the tag of a virtual or physical nature?
- For example, a virtual address or a physical address.
- Another example, a cache of data in a file or physical disk sectors.
- Different behavioral characteristics

# Page Coloring

# Deeper dive into page coloring



## Final comment

- “Caches work great, except when they don’t.”